

Pushing forward the dimension of fcc lattices

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(joint work with Jean-Marie Maillard)

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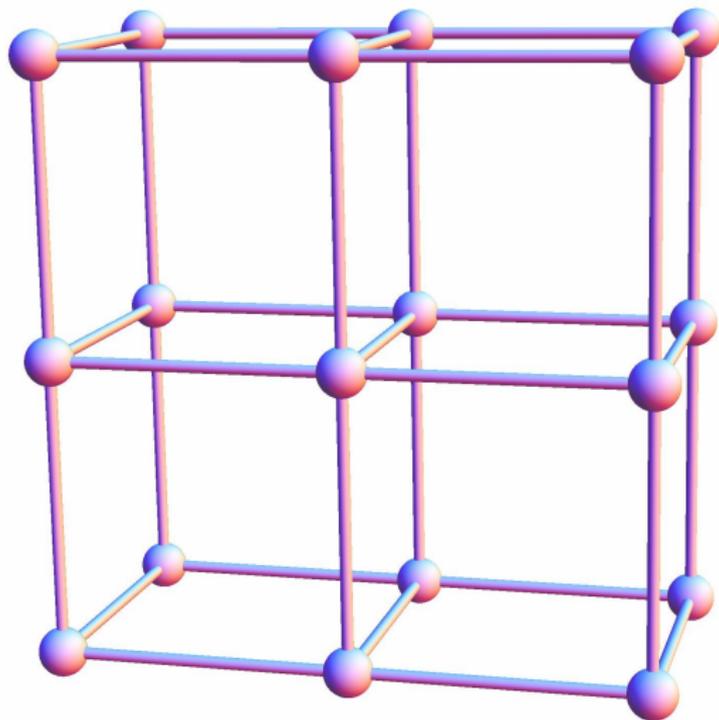


Face-Centered Cubic (fcc) Lattice

Motivation: statistical physics, crystallography, atomic structure

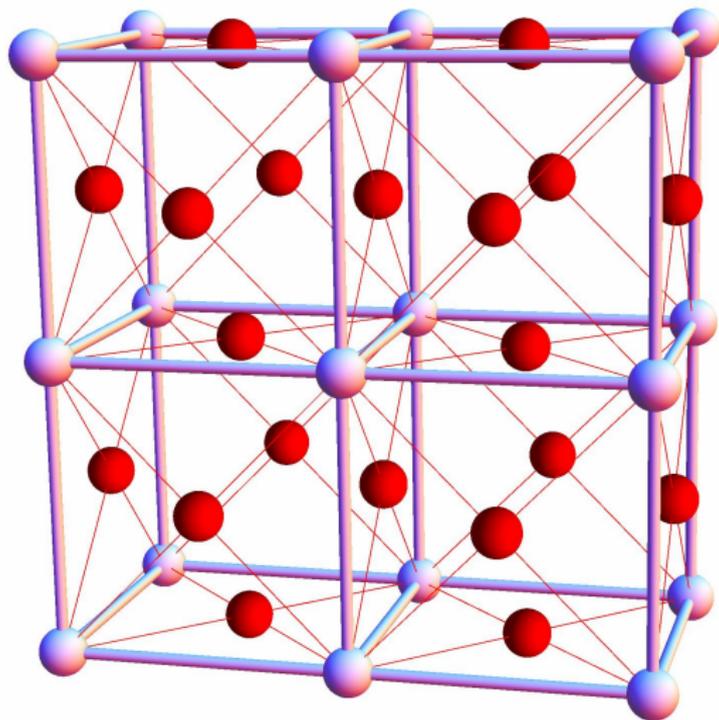
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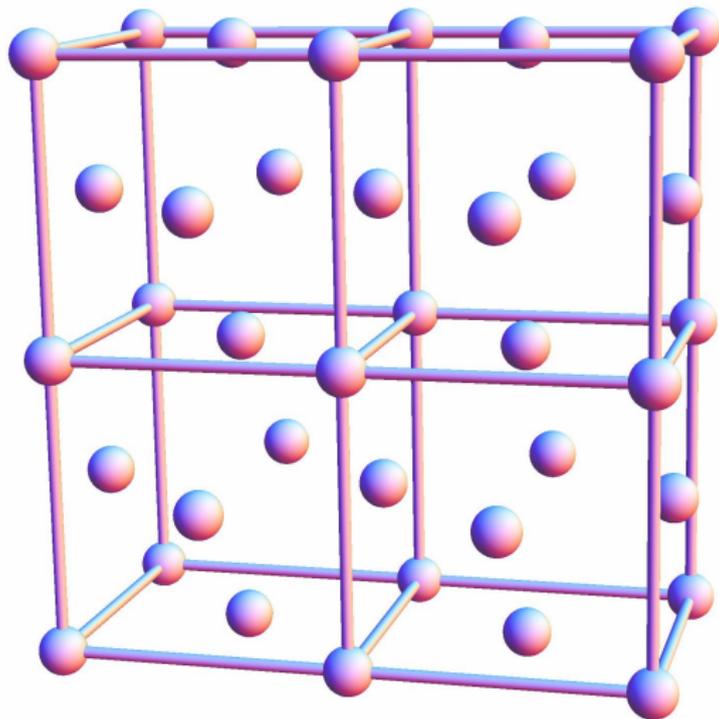
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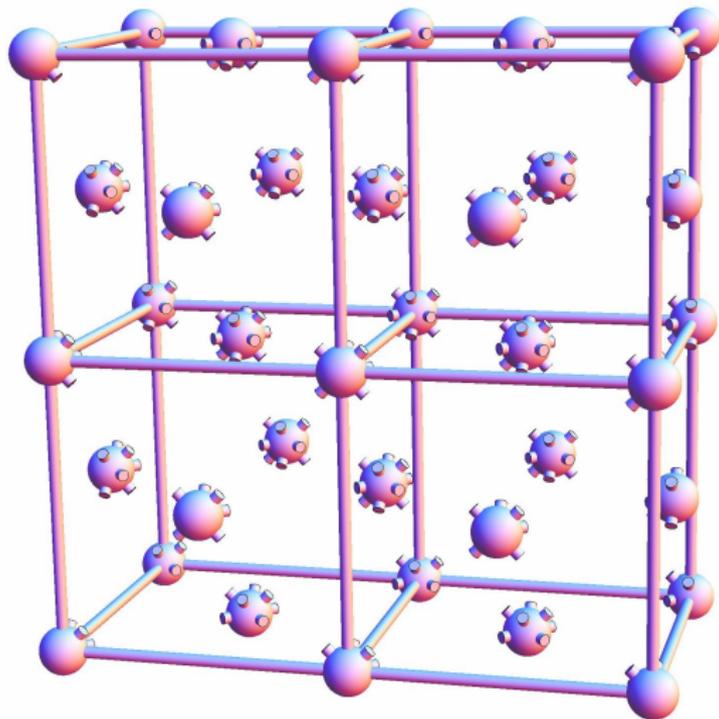
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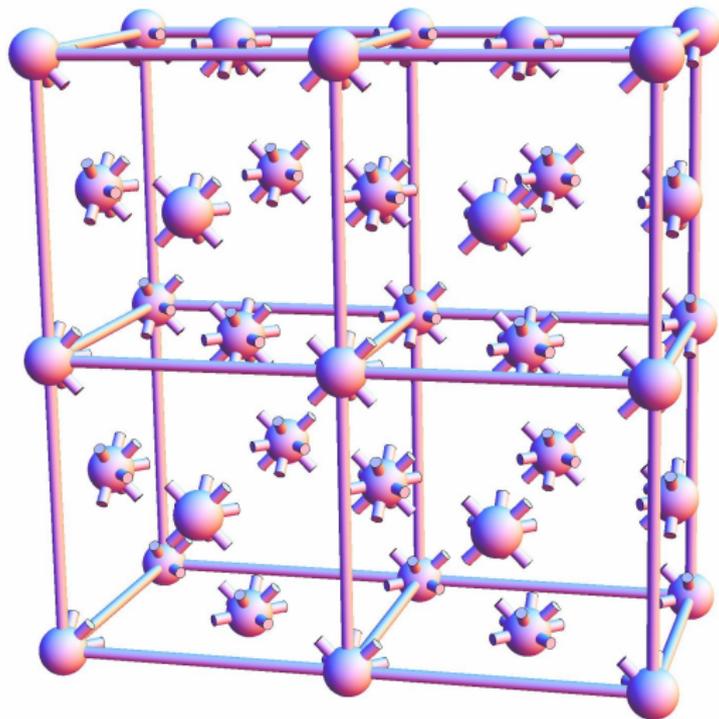
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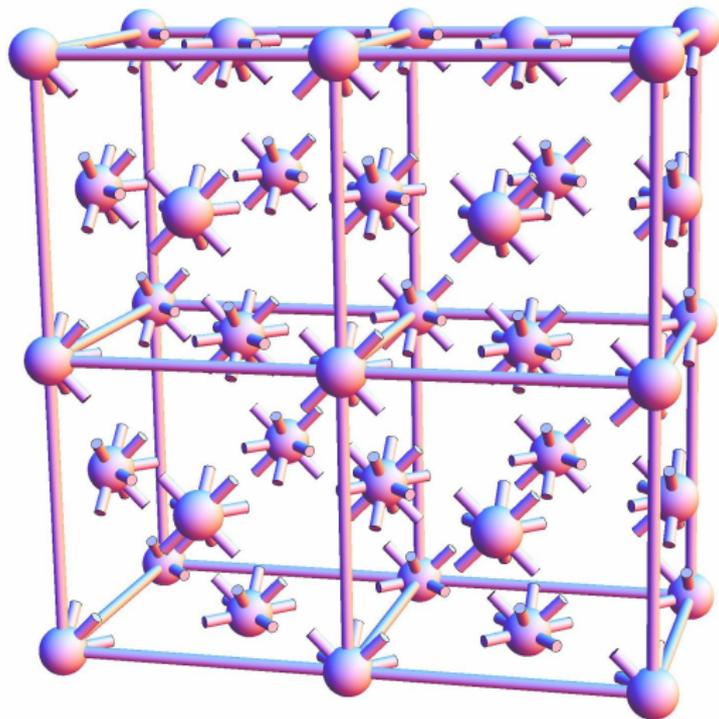
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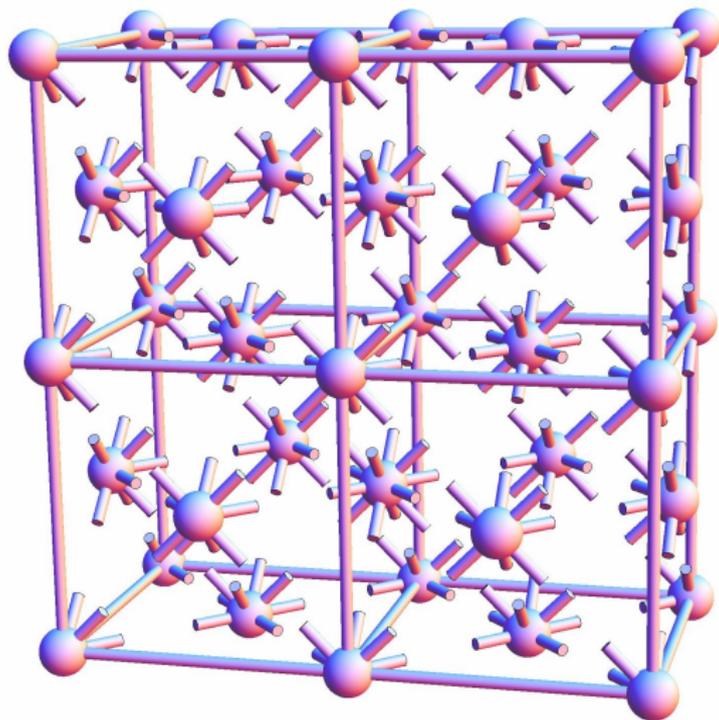
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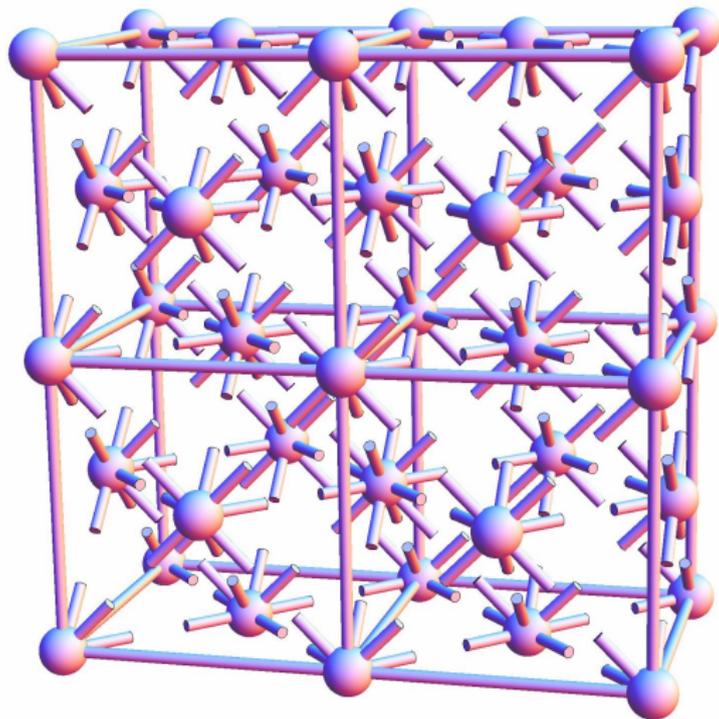
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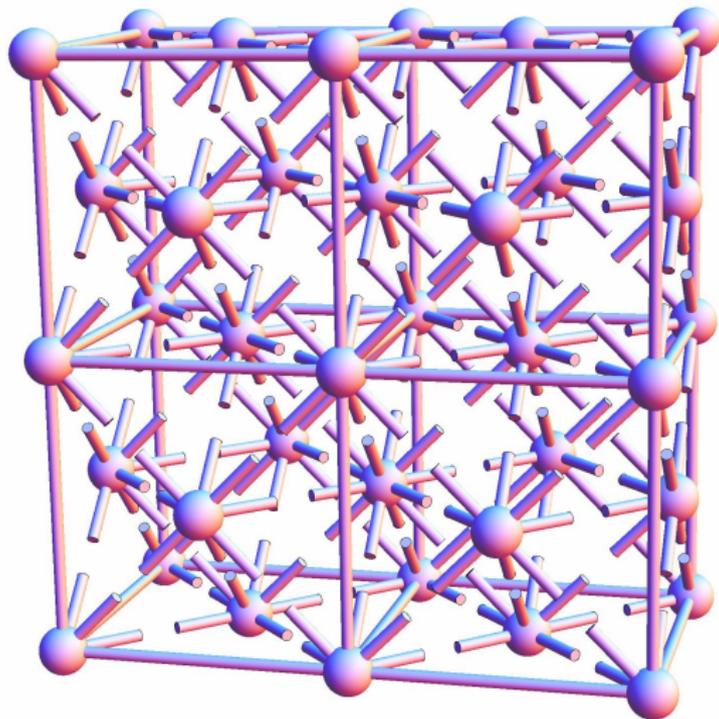
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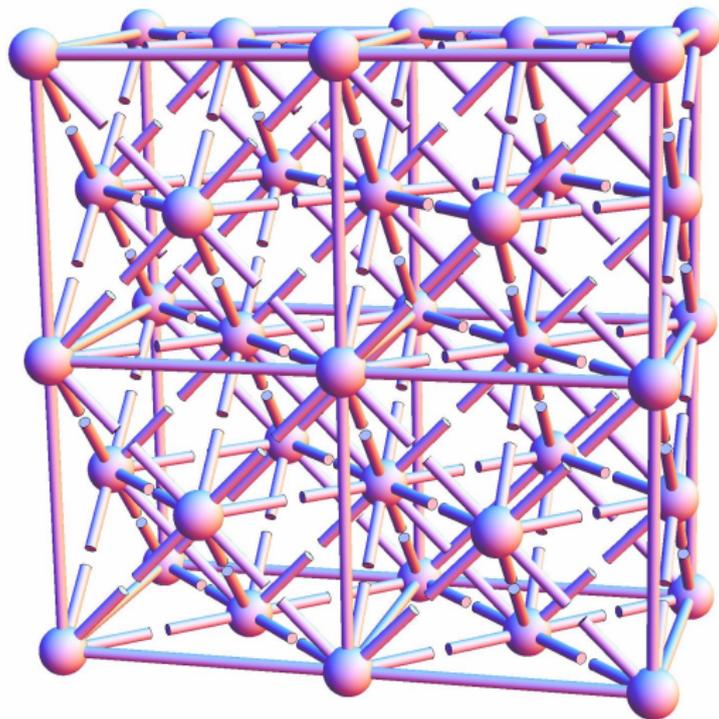
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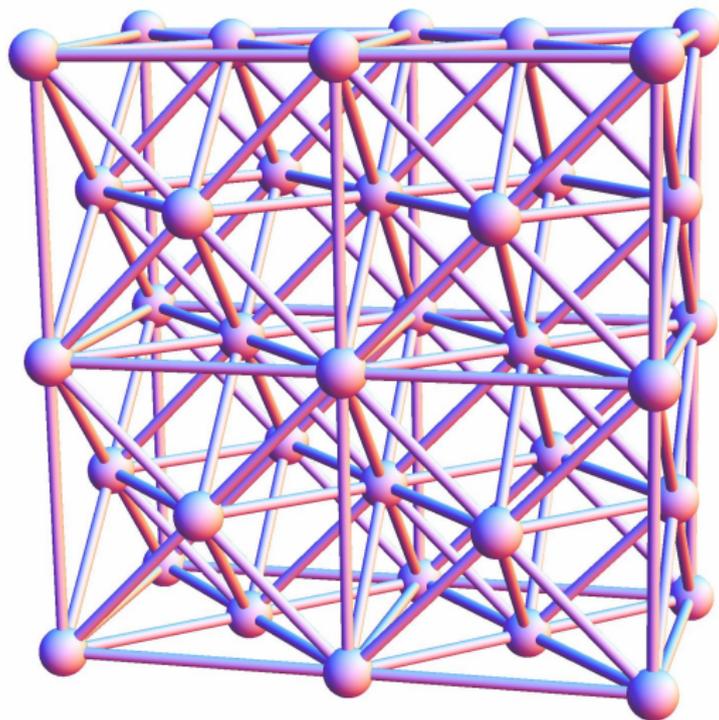
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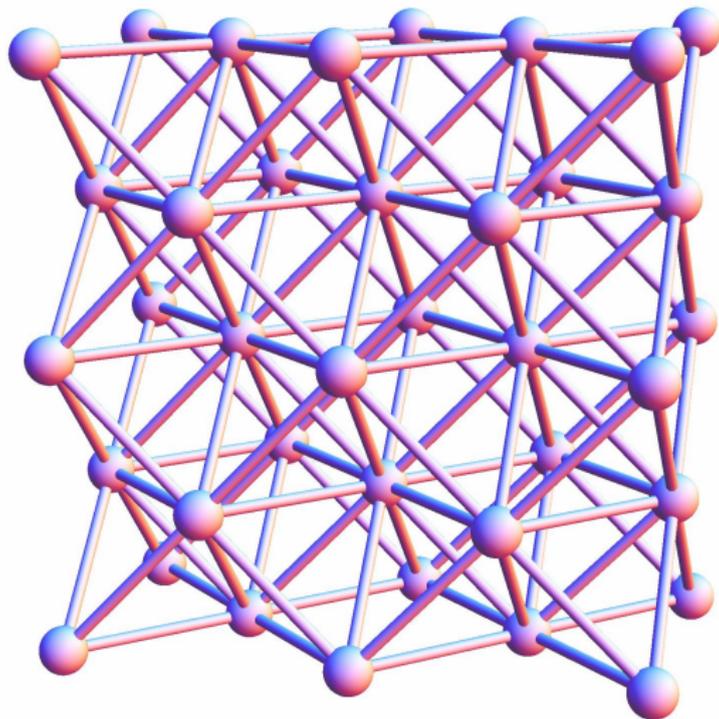
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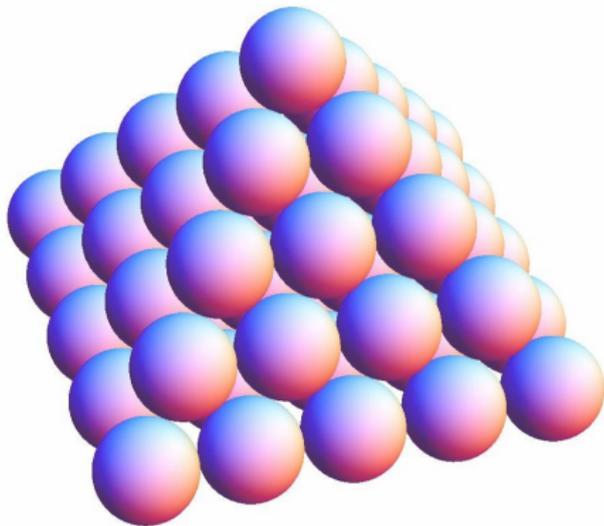
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Densest possible packing of spheres: Kepler conjecture
(proved by Hales in 2005)

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The d -dimensional fcc lattice is composed of $1 + \binom{d}{2}$ translated copies of \mathbb{Z}^d :

$$\mathbb{Z}^d + \sum_{1 \leq i < j \leq d} \left(\frac{1}{2}(\mathbf{e}_i + \mathbf{e}_j) + \mathbb{Z}^d \right).$$

Random Walks

We consider random walks on the lattice points:

- ▶ Move to one of the nearest neighbors in each step.
- ▶ All steps have the same probability.
- ▶ A lattice point can be visited arbitrarily often.
- ▶ Starting point is the origin $\mathbf{0}$.

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The set of permitted steps in the d -dimensional fcc lattice is

$$S = \{(s_1, \dots, s_d) \in \{0, -1, 1\}^d : |s_1| + \dots + |s_d| = 2\},$$

i.e., there are $4\binom{d}{2}$ steps (called the *coordination number* c).

Lattice Green's Function

The *lattice Green's function* is the probability generating function

$$P(\mathbf{x}; z) = \sum_{n=0}^{\infty} p_n(\mathbf{x}) z^n$$

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Its multivariate generating function

$$F(\mathbf{y}; z) = \sum_{n=0}^{\infty} \sum_{\mathbf{x} \in \mathbb{Z}^d} p_n(\mathbf{x}) \mathbf{y}^{\mathbf{x}} z^n = \frac{1}{1 - \frac{z}{|S|} \sum_{\mathbf{s} \in S} \mathbf{y}^{\mathbf{s}}}$$

is rational (S is the step set).

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We are particularly interested in

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that encodes the return probabilities. It is D-finite.

Here $\lambda(\mathbf{k})$ is called the *structure function* of the lattice; it is given by the discrete Fourier transform of the single-step probabilities:

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In the case of the fcc lattice:

$$\lambda(\mathbf{k}) = \binom{d}{2}^{-1} \sum_{1 \leq i < j \leq d} \cos(k_i) \cos(k_j).$$

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For $d = 3$, the return probability is one of *Watson's integrals*:

$$R = 1 - \left(\frac{1}{\pi^3} \int_0^\pi \int_0^\pi \int_0^\pi \frac{dk_1 dk_2 dk_3}{1 - \frac{1}{3}(c_1 c_2 + c_1 c_3 + c_2 c_3)} \right)^{-1} = 1 - \frac{16 \sqrt[3]{4} \pi^4}{9(\Gamma(\frac{1}{3}))^6}$$

where $c_i = \cos(k_i)$.

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The Taylor coefficients can be obtained in various ways:

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- ▶ Walk enumeration
- ▶ Reduction of dimension by multi-level guessing
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3. WZ proof theory (creative telescoping)

Previous results

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7	11	68	$\lambda^7(\lambda - 1)^3(\lambda - 2)$

Expansion of the integrand

Example for $d = 3$ (c_i denotes $\cos k_i$)

Expand the integrand in a geometric series:

$$\frac{1}{1 - \frac{z}{3}(c_1c_2 + c_1c_3 + c_2c_3)} = \sum_n \left(\frac{z}{3}\right)^n (c_1c_2 + c_1c_3 + c_2c_3)^n$$

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Use Wallis's integration formula:

$$\frac{1}{\pi} \int_0^\pi \cos^{2n} k \, dk = 4^{-n} \binom{2n}{n}$$

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The n -th Taylor coefficient can be computed by a recursive method.

The Recursive Method

Define

$$T_d(n, j) := \frac{4^{n+j}}{\pi^d} \int_0^\pi \cdots \int_0^\pi \left(\sum_{1 \leq i < j \leq d} \cos(k_i) \cos(k_j) \right)^n \left(\sum_{i=1}^d \cos(k_i) \right)^{2j} d\mathbf{k}.$$

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A straightforward calculation leads to the recurrence:

$$T_d(n, j) = \sum_{p=0}^n \sum_{q=q_1}^{q_2} \binom{n}{p} \binom{2j}{2q+p-n} \binom{2n+2j-2p-2q}{n+j-p-q} T_{d-1}(p, q)$$

with $q_1 = [(n-p+1)/2]$ and $q_2 = [(n-p+2j)/2]$.

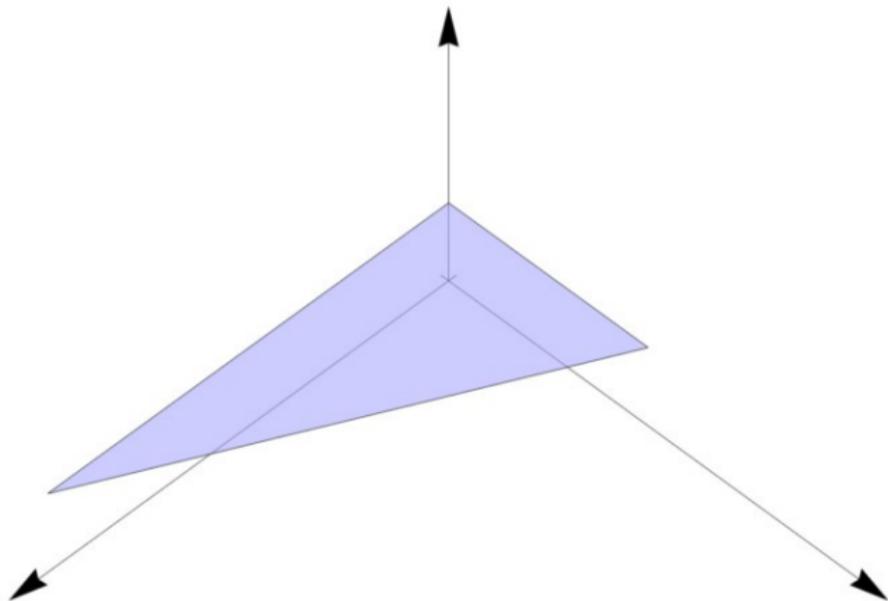
Initial condition:

$$T_2(n, j) = \sum_{p=p_1}^{p_2} \binom{2p}{p} \binom{2j}{2p-n} \binom{2n+2j-2p}{n+j-p}$$

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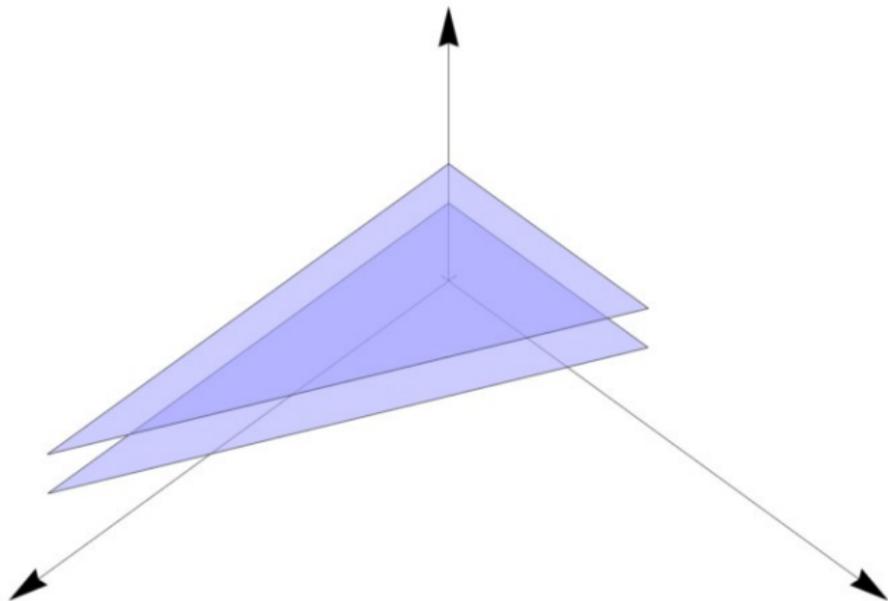
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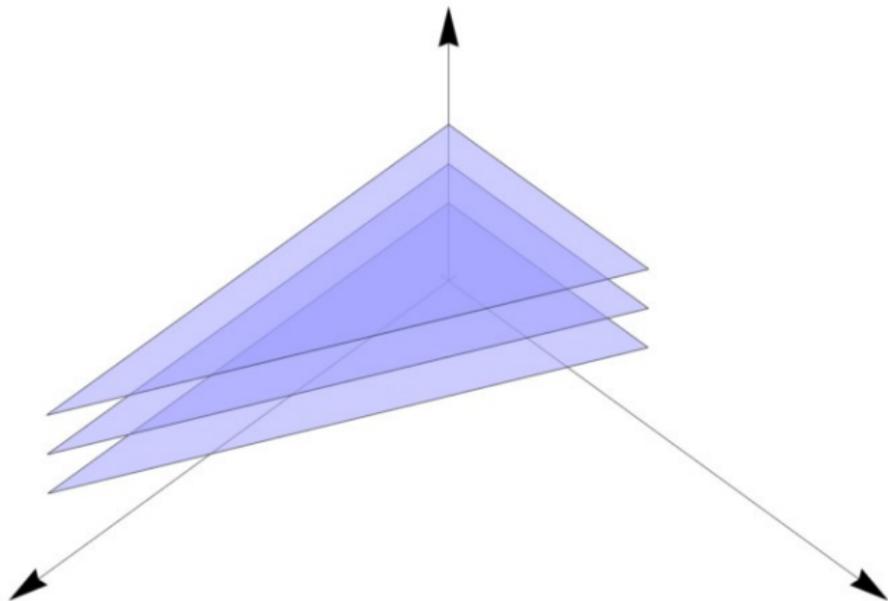
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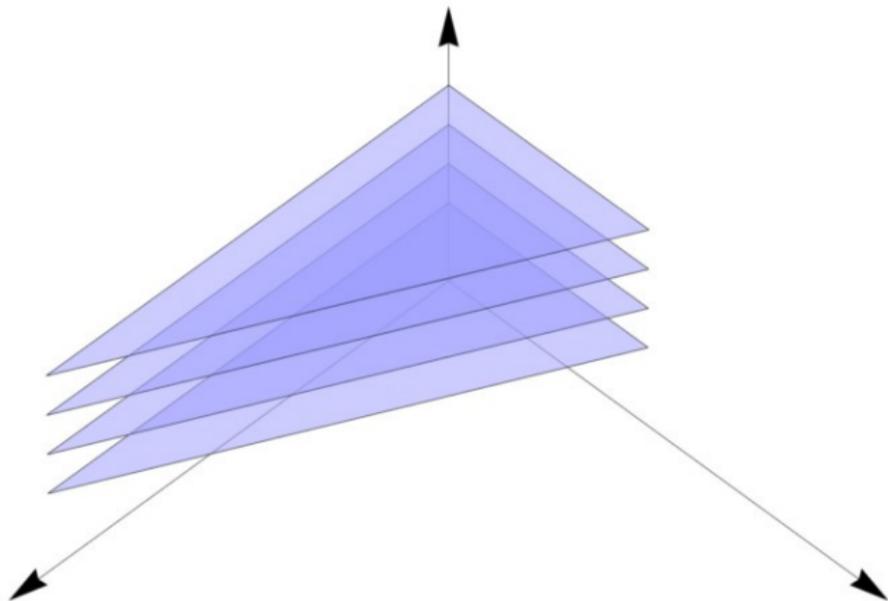
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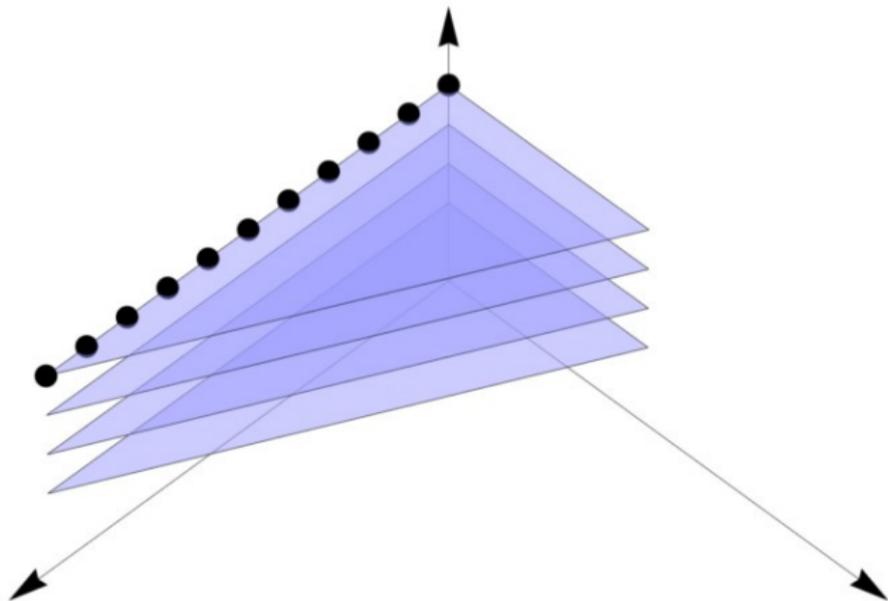
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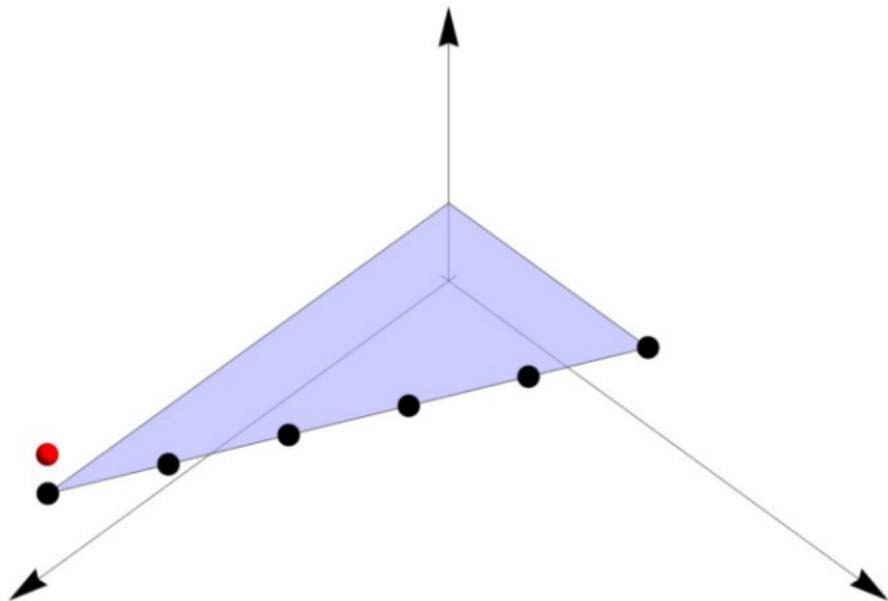
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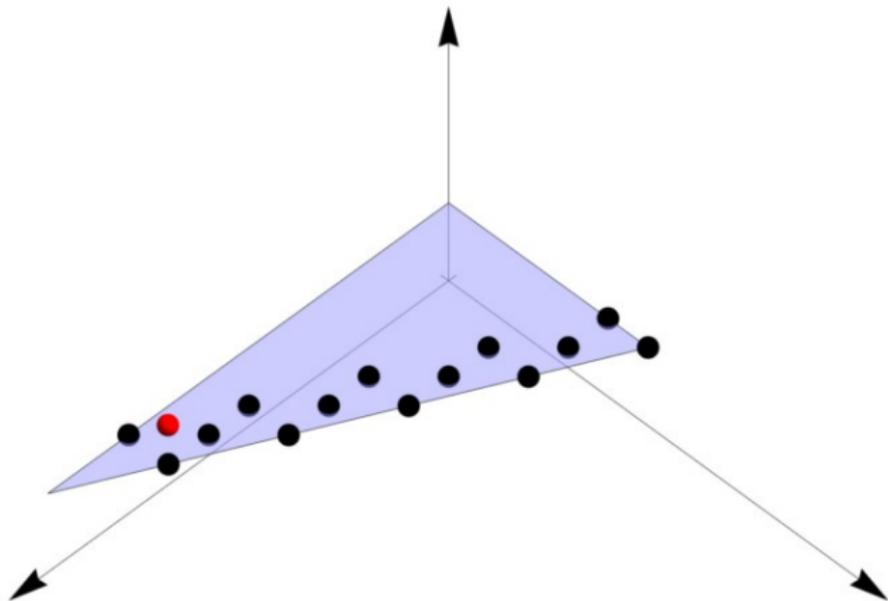
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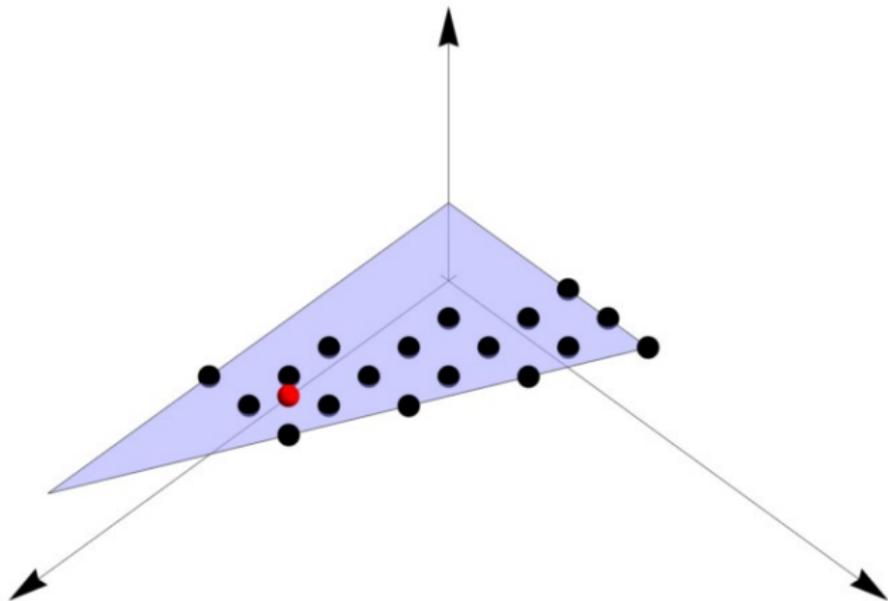
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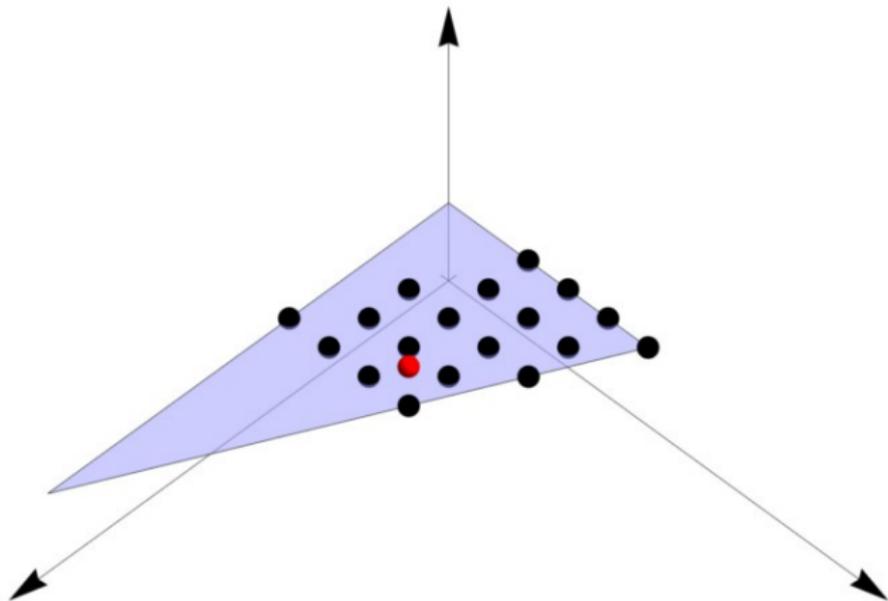
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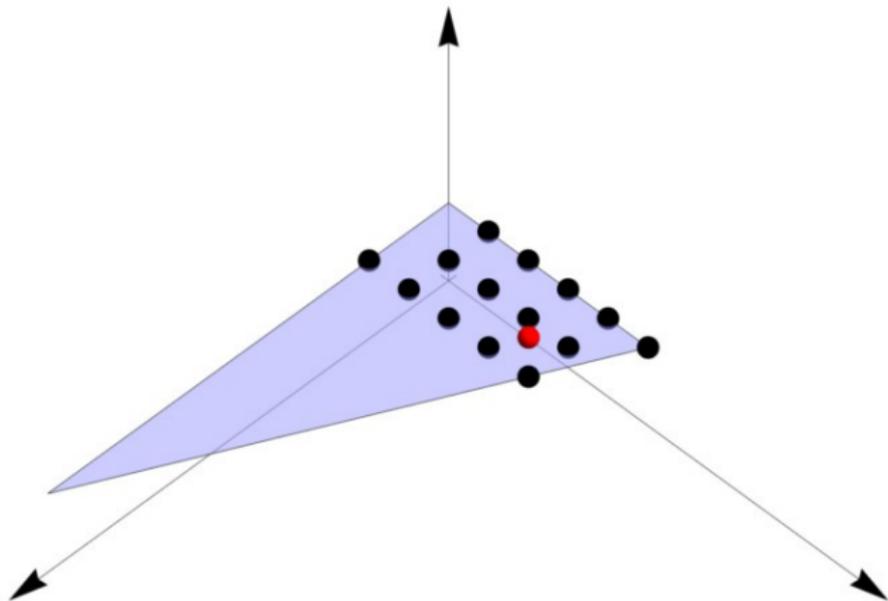
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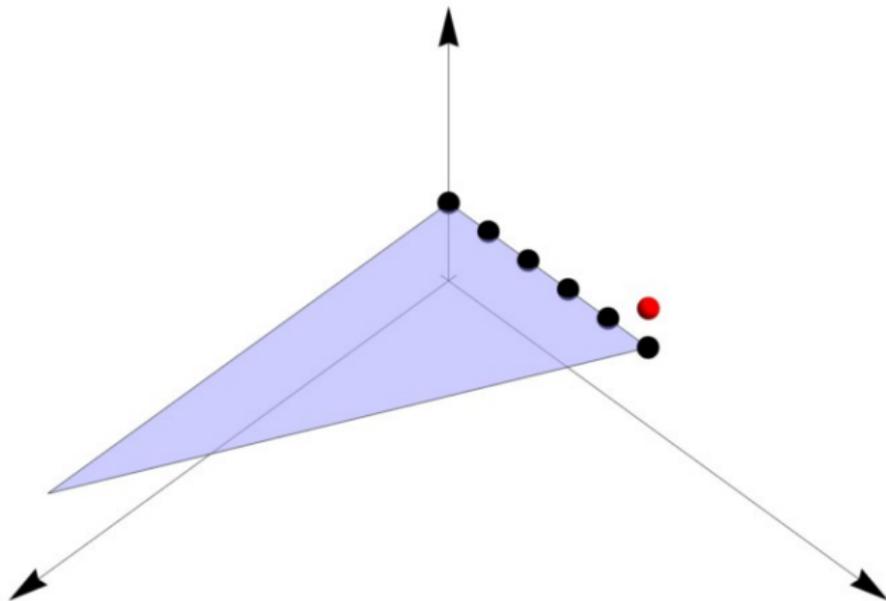
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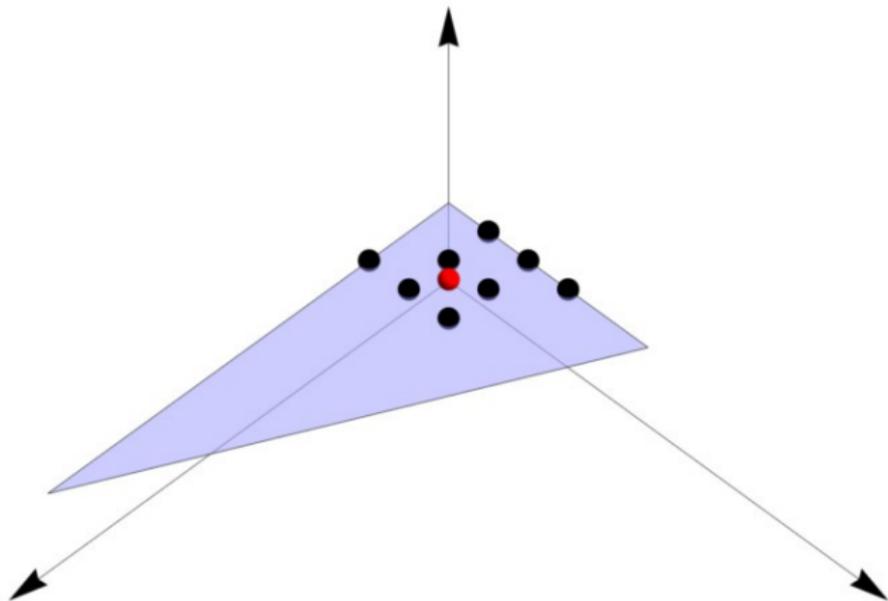
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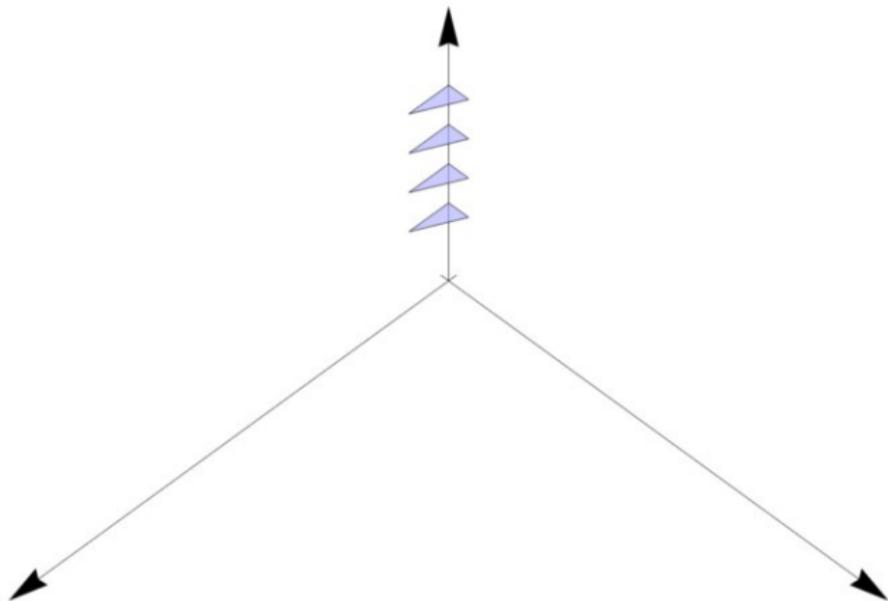
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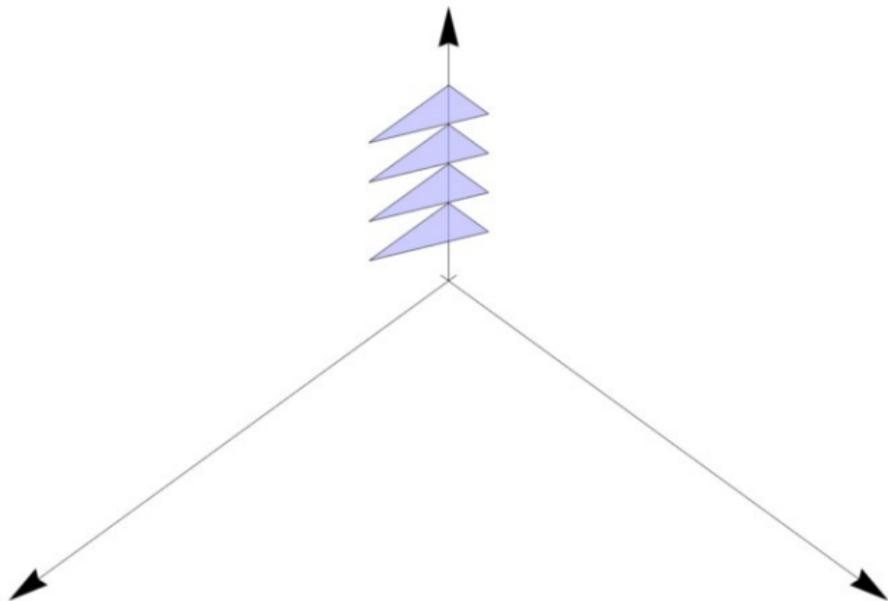
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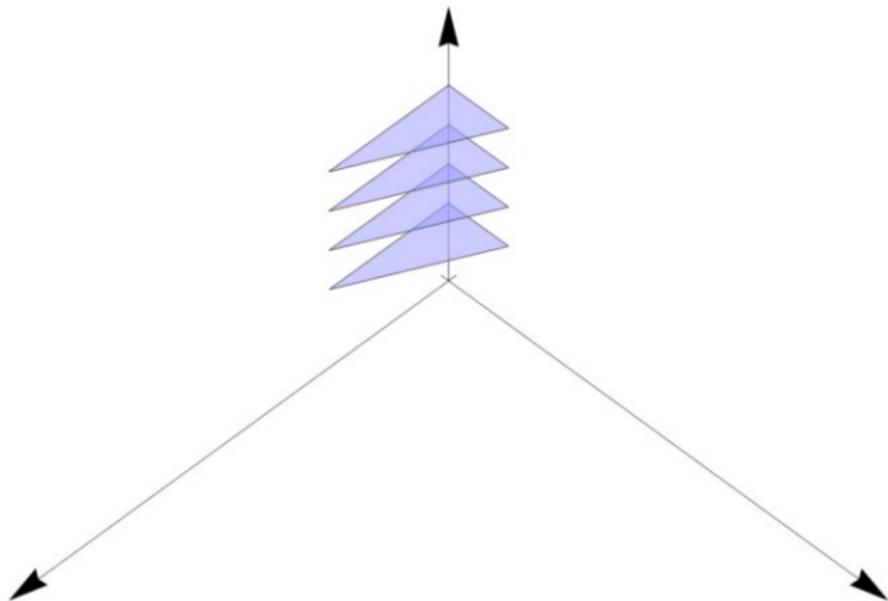
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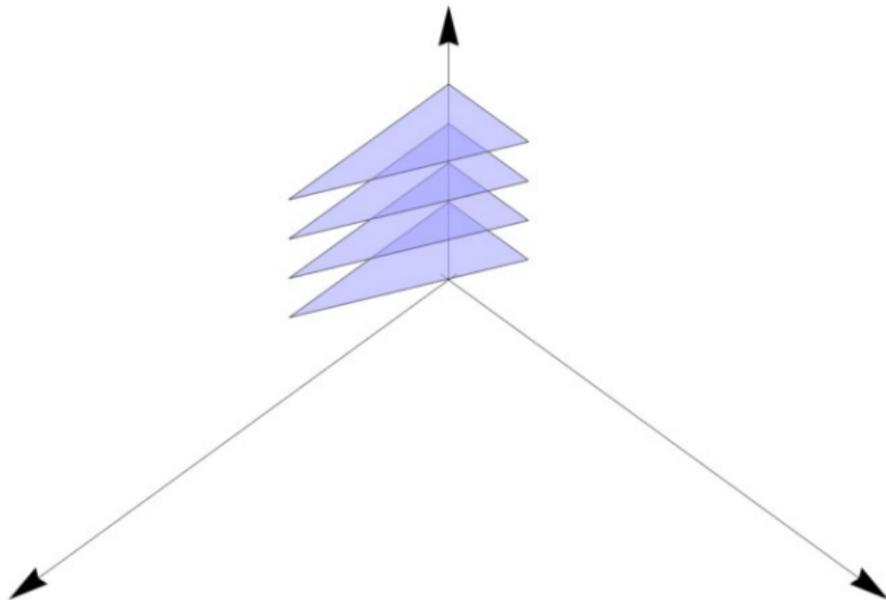
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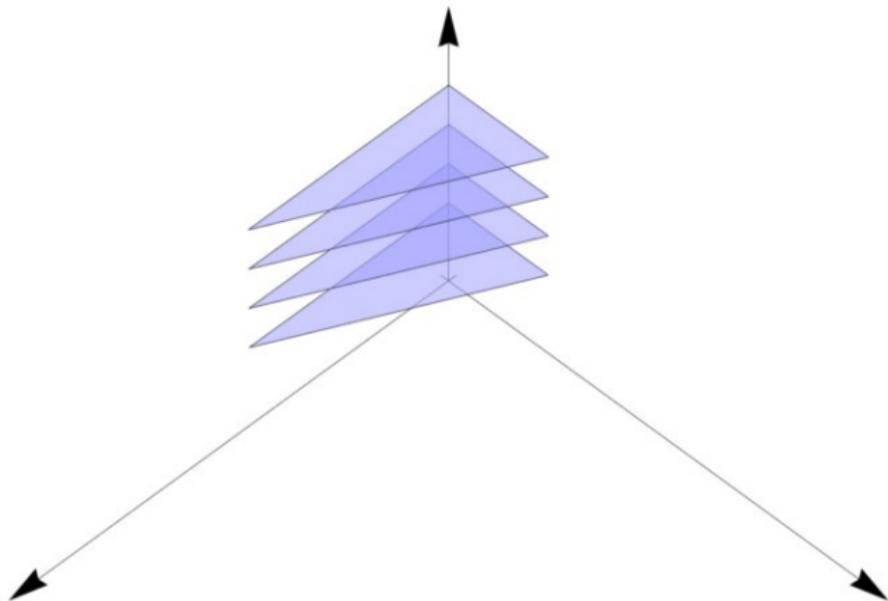
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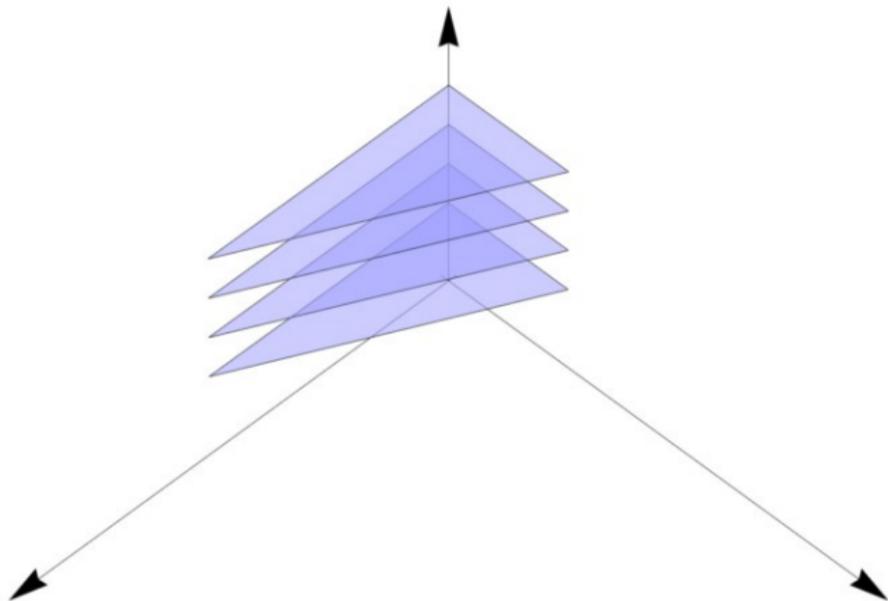
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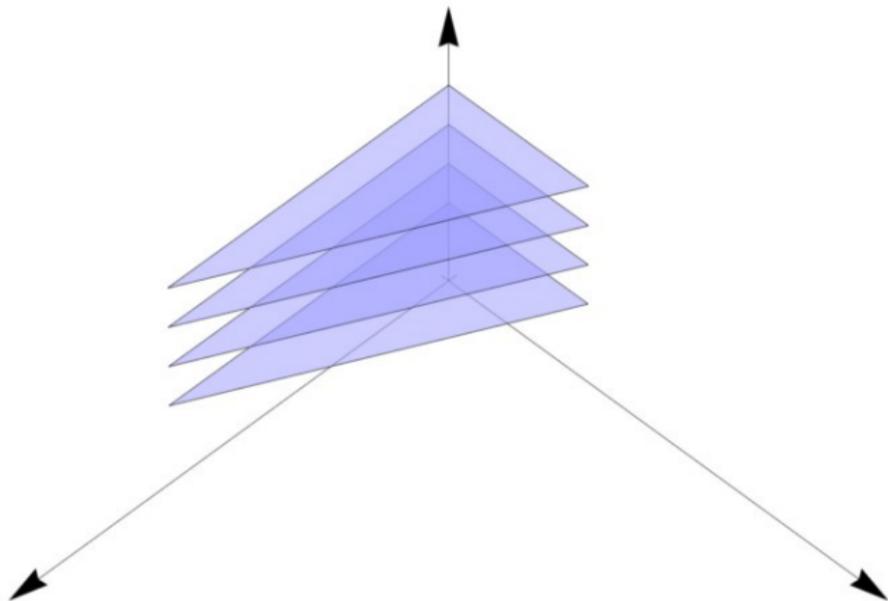
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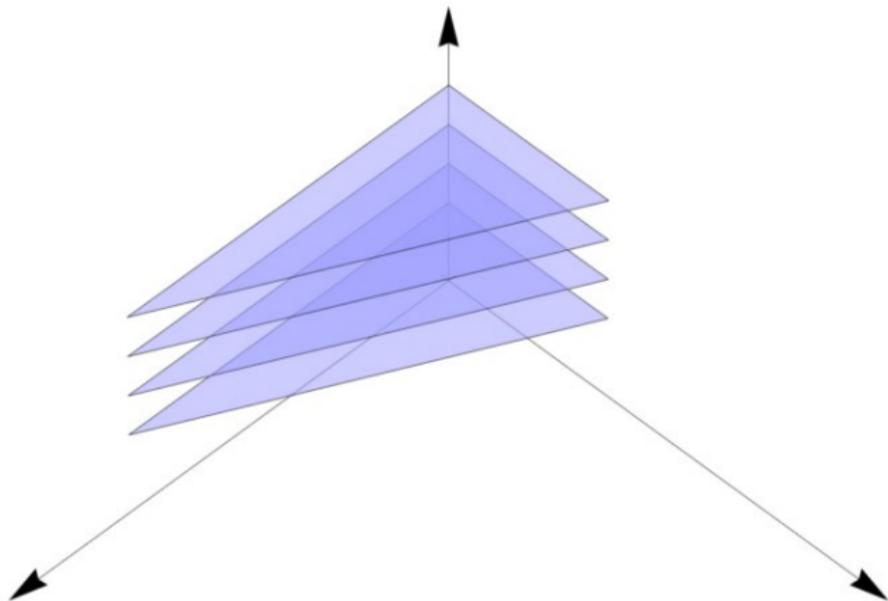
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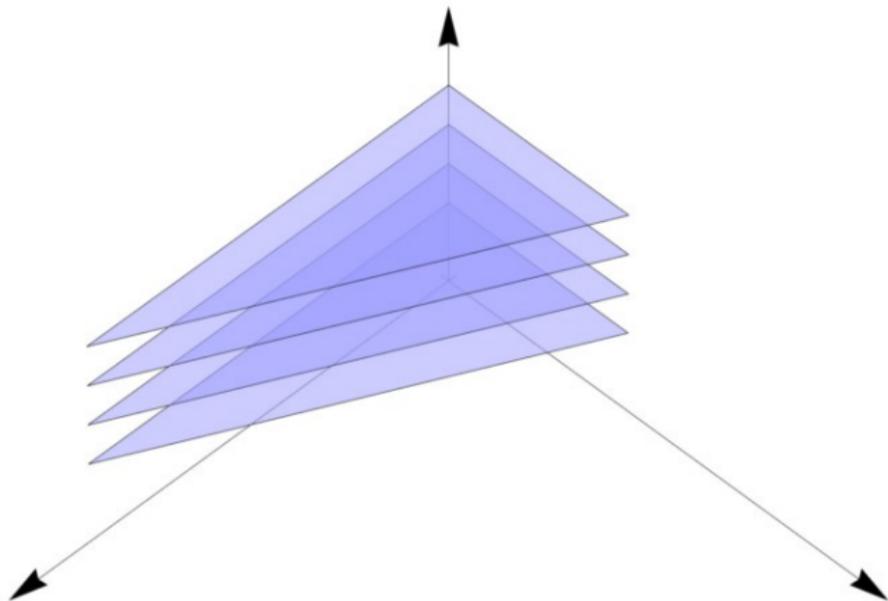
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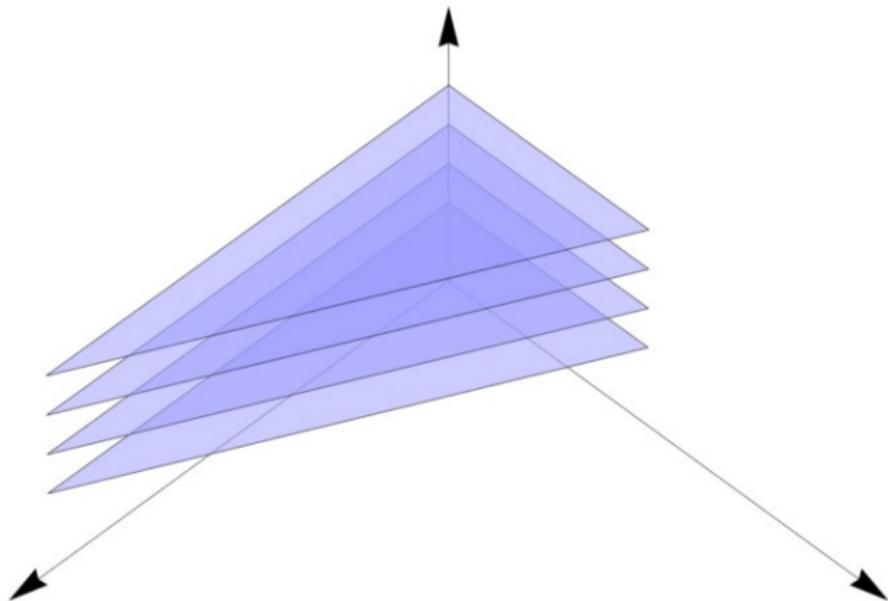
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1. Main loop over d :
 - ▶ Have to fix the number of Taylor coefficients $T_d(n, 0)$ a priori.
 - ▶ Recompute the coefficients in each iteration.
2. Main loop over n :
 - ▶ No knowledge about the necessary expansion order is needed.
 - ▶ Recomputation is avoided.
 - ▶ One needs to keep the full 3-dimensional array.

New Results

d	order	degree	terms	ind. pol. at $z = 0$
3	3	5	20	λ^3
4	4	10	40	λ^4
5	6	17	88	$\lambda^5(\lambda - 1)$
6	8	43	228	$\lambda^6(\lambda - 1)^2$
7	11	68	391	$\lambda^7(\lambda - 1)^3(\lambda - 2)$

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12*	32	617	3618	$\lambda^{12}(\lambda - 1)^8(\lambda - 2)^6(\lambda - 3)^4(\lambda - 4)^2$

* modulo prime

Landau Singularities

Leading coefficient of the order-6 operator for $d = 5$:

$$16(-5 + z)(-1 + z)z^4(5 + z)^2(10 + z)(15 + z)(5 + 3z) \\ \times (-675000 + 3465000z - 1053375z^2 + 933650z^3 \\ + 449735z^4 + 144776z^5 + 15678z^6)$$

Landau Singularities:

- ▶ Singularities of a function defined by a multiple integral
- ▶ Can be found by imposing conditions on the integrand.
- ▶ For $P(\mathbf{0}; z)$ one obtains

$$\binom{d}{2} \frac{2(1 - k)}{d^2 - (k + 4j + 1)d + 4j^2 + k + 4jk}$$

for $k = 0, 2, 3, \dots, d - 1$ and $j = 0, 1, \dots, [(d - 1)/2]$.

- ▶ Can be used as a consistency check.

Landau Singularities for $d = 11$

The leading coefficient of the differential operator is

$$\begin{aligned} & x^{22} (x + 11)^6 (55 + x)^2 (x - 1) (8x + 55) (29x + 55) \\ & (4x + 55) (2x + 55) (4x + 11) (7x + 165) (7x - 55) \\ & (2x + 33) (17x + 55) (x + 44) (13x + 275) (3x + 55) \\ & (7x - 11) (13x + 55) (7x + 110) (x + 35) (3x + 22) \\ & (x + 99) (19x - 55) (7x + 33) (9x + 11) (x + 15) \\ & (9x + 55) (17x + 275) (3x + 77) (23x + 165) \\ & \times \langle \text{irreducible polynomial of degree } 352 \rangle \end{aligned}$$

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3. For even dimension d :
 - ▶ Exterior square has a drop of order by 1.
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4. There exists a nontrivial homomorphism that maps the solutions of the operator L to the solutions of its adjoint:

$$\exists L_{\text{hom}} : \quad \text{adj}(L_{\text{hom}}) \cdot L = \text{adj}(L) \cdot L_{\text{hom}}$$

This gives rise to a “canonical decomposition”.

Canonical Decomposition

Perform successive Euclidean right divisions of L and L_{hom} :

$$L_0 := L$$

$$L_1 := L_{\text{hom}}$$

$$L_i := U_{i+1}L_{i+1} + L_{i+2}$$

- ▶ The quarks U_i are self-adjoint.
- ▶ Tower of intertwiners: $\text{adj}(L_{i+1}) \cdot L_i = \text{adj}(L_i) \cdot L_{i+1}$
- ▶ Also the L_i have special Galois groups.
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Example: ($d = 7, r = 11$)

$$L = (U_1U_2U_3U_4U_5 + U_1U_2U_5 + U_1U_4U_5 + U_1U_2U_3 + U_3U_4U_5 + U_5 + U_3 + U_1) \cdot r$$

where U_1, U_2, U_3, U_4 have order 1, the order of U_5 is 7, and r is a rational function.

Pushing forward the dimension of fcc lattices

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(joint work with Jean-Marie Maillard)

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